STOP/DJVU - Unpacking

Summary

In this post, I will walk through the unpacking process of a stop/djvu sample with the sample hash below:

SHA256 f02b45b579b65a1ea89f2d9443f2c1a1484dec0bc66591ff4d3ad6ce63d635aa

I got this sample from Malware Bazaar:

[https://bazaar.abuse.ch/sample](https://bazaar.abuse.ch/sample/f02b45b579b65a1ea89f2d9443f2c1a1484dec0bc66591ff4d3ad6ce63d635aa/) [/f02b45b579b65a1ea89f2d9443f2c1a1484dec0bc66591ff4d3ad6ce63d635aa/](https://bazaar.abuse.ch/sample/f02b45b579b65a1ea89f2d9443f2c1a1484dec0bc66591ff4d3ad6ce63d635aa/)

We will do:

- Static Analysis with Ghidra
- Dynamic Analysis with x32dbg, accompanied by static analysis in Cutter
- Dumping the payload with pe-sieve

This writeup was originally called "Teambot Unpacking", as Twitter user @gabbbarrr pointed out, this is, despite it's label, STOP/DJVU ransomware.

Static Analysis

After a quick look with DetectItEasy to check the entropy, it's easy to see this sample is packed:

So, I loaded the sample into Ghidra for a closer examination. Starting from the entry function, it isn't very difficult to find main. After calls to \Box setargy and \Box setenvp, a call to \Box wincmdln is made. Immediately after this, the main function is called. This function is was not labled as main, so I renamed it as you can see in the image below:

```
DAT_00500de0 = GetCommandLineA();
DAT_004c0080 = __crtGetEnvironmentStringsA();
iVar1 = __setargv();
if (iVar1 < 0) {
  __amsg_exit(8);
iVar1 = _{setenvp();if (iVar1 < 0) {
  __amsg_exit(9);
iVar1 = cinit(1);
if (iVar1 != 0) {
  __amsg_exit(iVar1);
  wincmdln();
local_24 = main();
```
The main function contains various API calls. However, these seem to be noise. Except for the two

functions before the return at the end, I found nothing significant here.

The first function, relabled by me to $\lfloor m \nu \rfloor$ load $\lfloor m \nu \rfloor$ loads $\lfloor m \nu \rfloor$ and $\lfloor m \nu \rfloor$ via LoadLibraryA:

The second function, labled as **FUN** 403906, is where the unpacking takes place.

After some initialization, **LocalAlloc** is retrieved from $msimq32$.dll via GetProcAddress. The resulting function address of LocalAlloc is stored in **addr localalloc**. With this, some memory is allocated, pointed to by **addr**. I labled the next function $\vert \mathbf{w} \vert$ virtual protect. This function changes the permissions of the allocated memory region to execute, read, write:

```
h_{\text{max}} = GetModuleHandleA((LPCSTR)&msimg32_dll);
_addr_localalloc = GetProcAddress(_h_msimg32,"LocalAlloc");
addr = (LPVOID)(*_addr\_localalloc)(0, size);mw_virtualprotect();
void mw_virtualprotect(void)
\{DWORD old_val;
 undefined4 PAGE_EXECUTE_READWRITE;
 PAGE_EXECUTE_READWRITE = 0x40;
                    /* PAGE_EXECUTE_READWRITE */
 VirtualProtect(addr,size,0x40,&old_val);
  return;
```
Continuing to analyze FUN 403906 reveals the function which unpacks the malware, thus, I renamed this to mw unpack :

In this function, the unpacking routine is applied. It begins with a large number of hard-coded 32-bit constants and ends with a loop containing the unpacking logic. The unpacked code is stored in the previously allocated memory.

```
local_1c0 = 0x5e24fb8f;local_11c = 0x13cc9676;local_128 = 0 \times 145ce368;local_c4 = 0x53a4e4aa;local_178 = 0 \times 180454cc;local_1b4 = 0x5abc23bd;local_74 = 0x7a98d7b;local_228 = 0 \times 48a0c18f;local_f0 = 0x6e4572f2;local_170 = 0x924e26a9;local_8 = (local_2d4 >> 5) + local_2ec ' local_2d8 ' local_2d0 + uVar2;if (size == 0x1a3) {
   MoveFileW((LPCWSTR)0x0,(LPCWSTR)0x0);
  ł
 uVar4 = uvar4 - local_8;local_10 = 4;uVar3 = uVar4 \times 0 \times 10 + local_2e4;local_8 = (uVar4 \gg ((byte)local_c 6 0x1f)) + local_2e8;uVar1 = local_2d0 + uVar4;if (size == 0xb3f) {
    GetConsoleAliasesA(local_6ec,0,(LPSTR)0x0);
    InterlockedPushEntrySList((PSLIST_HEADER)0x0,(PSINGLE_LIST_ENTRY)0x0);
 \mathcal{F}_DAT_004c0b98 = 0;uVar2 = uVar2 - (uVar3 \text{ '} uVar1 \text{ '} local_8);local_2d4 = uVar2;FUN_00402fc3();
 local_2dc = local_2dc + -1;} while (local_2dc != 0);
param_1[1] = uVar4;
\starparam_1 = uVar2;
return;
```
I didn't dig through the unpacking algorithm itself very much.

Eventually, before the return of FUN 403906, the unpacked code is executed in exec unpacked.

```
mw\_unpack();
iVar2 = 0;do \{GetLastError();
  if (iVar2 == 0x770e) {
    FUN_004038b5();
  J.
  iVar2 = iVar2 + 1;} while (iVar2 < 0x286b97d);
iVar2 = 0x7b;do \{if (size == 0 \times d) {
    CreateDirectoryA((LPCSTR)0x0,(LPSECURITY_ATTRIBUTES)0x0);
    lstrlenA("Yukanevakuriya duhifufacisubop");
    CloseEventLog((HANDLE)0x0);
  \mathcal{E}iVar2 = iVar2 + -1;} while (iVar2 != \emptyset);
exec\_unpackage();
return;
```
The exec unpacked was again labled as such by me:

What I pieced together above can be confirmed by debugging in x32dbg:

LocalAlloc - allocates memory at address 0xb50020

eax=00B50020

text:00404079 tb/exe:\$4079 #3479

Virtual Protect - applied to the memory area 0xb50020, set permissions to ERW

After unpacking, a jump to the unpacked payload occurs.

That's as far as I went with the static analysis, next, let's fire up x32dbg and get that payload out.

Dynamic Analysis

I will restart the debugging process here, the debugging in the "Static Analysis" section was only to link the static analysis findings with the actual execution flow of the program.

After I restarted the debugging process, I set a breakpoint at *VirtualProtect*. That way, I can get the address of the allocated memory and break before the unpacked code is executed.

As you can see from the stack, the permissions are set to execute, read, write (0x40) and the allocated memory's address is the topmost argument.

I set a hardware breakpoint at the memory address in the dump.

If you hit continue now, you will land in the routine before the part that executes the unpacked code:

We're looking for the call at 0x404139, this calls the unpacked code, remember FUN 403906 from the static analysis section, this is the routine we are in, so we can find 0x404139 by scrolling down. I set a breakpoint here and removed the hardware breakpoint. After that, I hit continue and ended up at the call at 0x404139.

After stepping into this call, we want to follow the execution from $\lim_{n \to \infty} \frac{1}{n}$

After following this, I landed in the allocated memory, where the unpacked code resides.

From here, there are two possible ways to continue.

- 1. Step through the code until we find something interesting
- 2. dump the memory of the shellcode and check it out in a disassembler.

I prefer Option 1.

Examine the Shellcode - 1

Dumping the memory region where the shellcode is stored and loading the dump into Cutter reveals several interesting routines. First this routine, which looks complicated judging by it's graph overview:

From my examination, it seems this performs further unpacking. This is confirmed later.

Also, there is a function containing stack strings:

```
mov dword [rbp + rax - 0x30], 0x6e72656b; 'kern'
mov eax, dword [var_38h]
add eax, 4
mov dword [var_38h], eax
mov eax, dword [var_38h]
mov dword [rbp + rax - 0x30], 0x32336c65; 'el32'
mov eax, dword [var_38h]
add eax, 4
mov dword [var_38h], eax
mov eax, dword [var_38h]
mov dword [rbp + rax - 0x30], 0x6c6c642e; '.dll'
```

```
mov dword [rbp + rax - 0x30], 0x74726956; 'Virt'
mov eax, dword [var_38h]
add eax. 4
mov dword [var_38h], eax
mov eax, dword [var_38h]
mov dword [rbp + rax - 0x30], 0x416c6175; 'ualA'
mov eax, dword [var_38h]
add eax, 4
mov dword [var_38h], eax
mov eax, dword [var_38h]
mov dword [rbp + ra \times - \theta x 30], \theta x 636f6c6c; 'lloc'
```
Continuing in the debugger, I put a breakpoint at VirtualAlloc and hit continue. Indeed, VirtualAlloc is called and some memory is allocated. I put a hardware breakpoint there and kept on stepping.

Once the hardware breakpoint was hit, I found myself in the unpacking routine shown in the graph overview above. I set a breakpoint at the return of this routine and watched the allocated memory being populated. Once the memory area was filled, I dumped that memory as well. As with the dump before, I loaded it into Cutter to examine it the code in detail.

Examine the Shellcode - 2

Looking at the strings of this dump shows that it contains the DOS header and DLL names among other interesting things.

0x001190d3 :: \$:(:4:8:<:@:D:H:L:P:T:X:\\:`:d:h:l:p:t:x:|: 0x000015ed !This program cannot be run in DOS mode.\r\r\n\$ 0x000cff08 R6017\r\n- unexpected multithread lock error\r\n

0x000d0b34 SetDefaultDIIDirectories 0x000f6b60 ADVAPI32.DLL 0x000f6b70 KERNEL32.DLL 0x000f6b80 NETAPI32.DLL 0x00100214 \\shell32.dll 0x00100314 kernel32.dll 0x001090c4 KERNEL32.dll 0x00109338 ADVAPI32.dll 0x0010942c OLEAUT32.dll 0x0010944c IPHLPAPI.DLL 0x00100198 Shell32.dll 0x00108bd2 WININET.dll 0x00108c76 SHLWAPI.dll 0x001093c8 SHELL32.dll 0x001094a2 CRYPT32.dll 0x000d06ac USER32.DLL 0x000d42d8 VAPI32.DLL 0x000f6c94 USER32.DLL 0x00108b06 RPCRT4.dll 0x001091de USER32.dll 0x0010945a WS2_32.dll 0x0010947e DNSAPI.dll 0x000cfc78 coree.dll 0x00100368 Psapi.dll 0x00108bec WINMM.dll 0x00109422 ole32.dll 0x00109a9e GDI32.dll 0x00108b46 MPR.dll 0x000fecf8 %s.dll

Looking through the functions in Cutter, one of them stands out. It contains stack strings of various APIs, among them, in this order:

- CreateProcessA
- GetThreadContext
- VirtualAlloc
- VirtualAllocEx
- VirtualFree
- ReadProcessMemory
- WriteProcessMemory
- SetThreadContext
- ResumeThread

call qword [var_dch] mov dword [var_b4h], eax

mov byte [var_280h], 0x43

mov byte [var_27fh], 0x72

mov byte [var_27fh], 0x72

; 'r'

mov byte [var_27ch], 0x65

; 'e'

mov byte [var_27ch], 0x74

; 't'

mov byte [var_27ch], 0x65

; 'e'

mov by mov dword [var_b4h], eax ; a
; 't'
; 'e'
; fcn.00000050
; 'r'
; 'c'
; 'e'
; 'e' $\frac{1}{2}$, $\frac{1}{3}$, $\frac{1}{3}$, $\frac{1}{4}$, $\frac{1}{3}$ mov byte [var_274h], 0x73 mov byte [var_273h], 0x41 mov byte [var_272h], 0 lea ecx, [var_280h] push rcx mov edx, dword [var_c4h] push rdx call qword [var_dch] mov dword [var_b0h], eax mov dword Lvar_boni, eax

mov byte [var_f4h], 0x47 ; 'G'

mov byte [var_f3h], 0x65 ; 'e'

mov byte [var_f2h], 0x74 ; 't'

mov byte [var_f1h], 0x54 ; 'T'

mov byte [var_f0h], 0x68 ; 'h'

mov byte [var_efh], 0x72 ; 'r'

mov $\frac{1}{2}$ $\frac{1}{2}$ $\frac{1}{2}$ $\frac{1}{2}$ mov byte [var_f4h], 0x47 $\frac{1}{2}$, $\frac{1}{2}$, $\frac{1}{2}$, $\frac{1}{2}$, $\frac{1}{2}$, $\frac{1}{2}$, $\frac{1}{2}$ mov byte [var_e6h], 0x78 mov byte [var_e5h], 0x74 mov byte [var_e4h], 0 lea eax, [var_f4h] وسامتها والمستحدث call gword [var_dch] mov dword [var_94h], eax $\begin{array}{cc}\n & \vdots & 'R' \\
 & \vdots & 'e' \\
 & \vdots & 's' \\
 & \vdots & 'u' \\
 & \vdots & 'm' \\
 & \vdots & 'e' \\
 & \vdots & \vdots & \vdots & \vdots & \vdots \\
 & \vdots & \vdots & \vdots & \vdots \\$ mov dword Evan _34n], eax
mov byte [van _25ch], 0x52
mov byte [van _25ah], 0x65
mov byte [van _25ah], 0x73
mov byte [van _259h], 0x75 mov byte [var_258h], 0x6d $\frac{1}{2}$ mov byte [var_2501], 0x00
mov byte [var_257h], 0x65 -) (T) $\frac{1}{2}$ $\frac{1}{2}$ $\frac{1}{2}$ $\frac{1}{2}$ in the $\frac{1}{2}$ $\frac{1}{2}$ $\frac{1}{2}$ e^{i} mov byte [var_253h], 0x65 $\begin{array}{c} 2 \overline{1} \\ 2 \overline{1} \\ 3 \overline{1} \end{array}$ mov byte [var_252h], 0x61 mov byte [var_251h], 0x64 mov byte [var_250h], 0 lea edx, [var_25ch]

This looks very much like Process Hollowing, which prompted me to set a breakpoint at ResumeThread.

I continued execution, hitting the breakpoint at ResumeThread , after which I opened Process Hacker:

The PID 5384, child of the original executable, is what was created by CreateProcessA. After the processes memory was written it should contain the stop/djvu payload. As ResumeThread is the last call for the Process Hollowing dumping PID 5384 with pe-sieve yields the unpacked stop/djvu malware.

I loaded the dumped executable into DetectItEasy, the entropy looks better:

... and there are several interesting imports:

Now on to analyzing the unpacked sample \sim o \sim - updates to follow ;-)

0xca7